

Introduction to Quantum Groups III

Reminder

Def (Rigid monoidal cat) A monoidal cat (C, \otimes, I) is "rigid" if all objects have left + right duals. ↙ aka "autonomous"

Def (Ribbon cat) A "ribbon cat" is a rigid braided monoidal cat (C, \otimes, I, τ) equipped w/ a "twist", ie. a nat trans $\theta: id_C \Rightarrow id_C$ w/ components $\theta_V: V \xrightarrow{\sim} V$ satisfying

- ① $\theta_{V \otimes W} = \tau_{WV} \tau_{VW} (\theta_V \otimes \theta_W)$
 - ② $\theta_I = id$
 - ③ $\theta_{V^*} = (\theta_V)^*$
- } "balancing axioms"

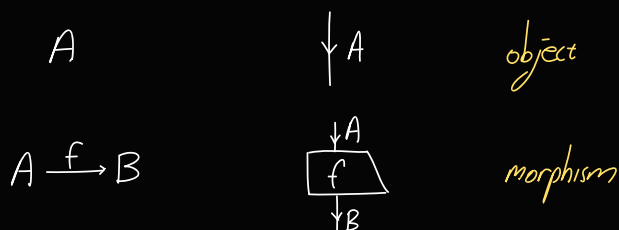
Lem A twist θ on a rigid braided monoidal cat (C, \otimes, I, τ) yields a nat trans $\delta: id_C \Rightarrow (-)**$ w/ components $\delta_V: V \xrightarrow{\sim} V^{**}$ s.t.

- ① $\delta_{V \otimes W} = \delta_V \otimes \delta_W$
 - ② $\delta_I = id$
 - ③ $\delta_{V^*} = (\delta_V^*)^{-1}$
- + vice-versa. ↙ ie. δ and θ are equivalent data
- δ is called a "pivotal structure"

§7. From Categories to Links

Lightning Tour (Graphical languages) † Selinger (2009)

① Category



$f \circ g$



vertical comp

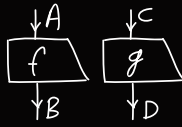
② Monoidal Category

$A \otimes B$



} horizontal comp

$f \otimes g$



I

(empty)

unit object

$A \otimes C \xrightarrow{f} B \otimes D$



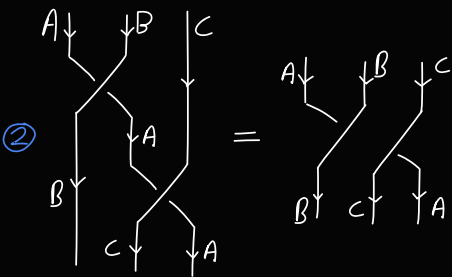
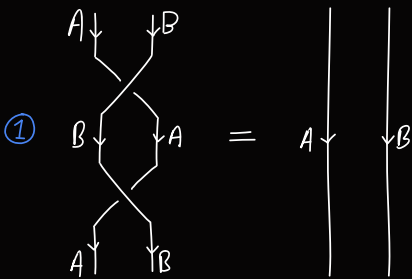
③ Braided Monoidal Category

$A \otimes B \xrightarrow{\tau_{AB}} B \otimes A$

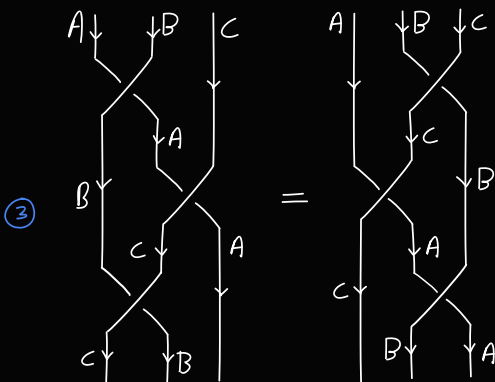


braiding

Some equations




hexagon axiom




Yang-Baxter

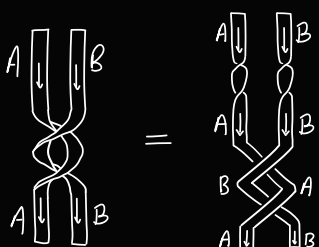
④ Balanced Braided Monoidal Category

A  object

$A \xrightarrow{\theta_A} A$  twist

$A \otimes B \xrightarrow{\theta_{A \otimes B}} B \otimes A$ 

Some equations

① 

balancing axiom
 $\theta_{v \otimes w} = \tau_{wv} \tau_{vw} (\theta_v \otimes \theta_w)$

⑤ Rigid Monoidal Category

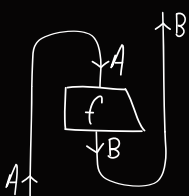
$A^*, *A$  duals

$I \xrightarrow{i_A} A^* \otimes A$  caps

$I \xrightarrow{i_A} A \otimes *A$  caps

$A \otimes A^* \xrightarrow{e_A} I$  cups

$*A \otimes A \xrightarrow{e_A} I$  cups

f^*  adjoints

Some equations

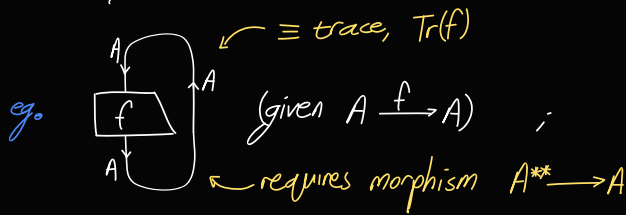
$$\textcircled{1} \quad \begin{array}{c} \downarrow A \\ \downarrow A \\ \downarrow A \end{array} \equiv \begin{array}{c} \downarrow A \\ \downarrow A \\ \downarrow A \end{array} \equiv \begin{array}{c} \downarrow A \\ \downarrow A \\ \downarrow A \end{array}$$

yanking identities

$$\textcircled{2} \quad \begin{array}{c} \uparrow A \\ \uparrow A \\ \uparrow A \end{array} \equiv \begin{array}{c} \uparrow A \\ \uparrow A \\ \uparrow A \end{array} \equiv \begin{array}{c} \uparrow A \\ \uparrow A \\ \uparrow A \end{array}$$

a remark

some apparently-innocent diagrams are not actually well-formed,



this motivates:

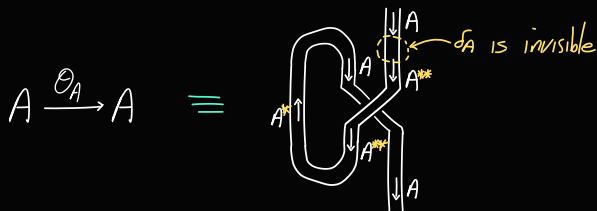
⑥ Pivotal Rigid Monoidal Category

$$A \xrightarrow{\delta_A} A^{**} \quad \downarrow A$$

a remark

we have seen that, in a rigid braided monoidal cat, twist $\Theta \leftrightarrow$ pivotal structure δ ;

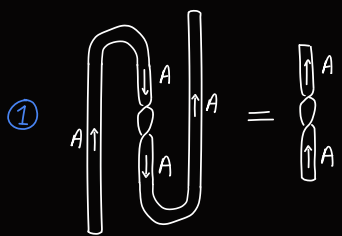
eg. given δ , define Θ by



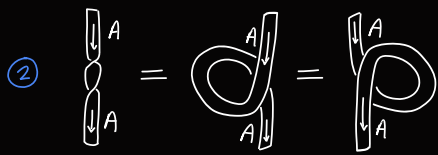
⑦ Pivotal/Balanced Rigid Braided Monoidal Category \equiv "Ribbon cat"

aka "tortile" cat

Some equations



balancing axiom
 $\theta_{A^*} = (\theta_A)^*$



relationship %
 θ & δ

Thm[?] (Coherence for ribbon cats) An equation relating ribbon cat morphisms follows from the axioms iff it holds in the graphical calculus, up to 3D framed isotopy

↑ special case proven in Shum (1994); I can't find general case proof

Convention Hereafter, adopt blackboard framing, + draw framed links single-stranded

§8. From Hopf Algebras to Links ← Kassel (1995), ch 13-14, Geiss (2014) §8.1.4

Def (Ribbon Hopf alg) A "ribbon Hopf alg" is a braided Hopf alg

$(A, m, \eta, \Delta, \varepsilon, S, R)$ w/ some $\theta \in Z(A)$ s.t.

① $\theta^2 = u S(u)$

② $\Delta(\theta) = (R_2 R)^{-1} [\theta \otimes \theta]$

③ $S(\theta) = \theta$

④ $\varepsilon(\theta) = 1$

} $u S(u) \in Z(A)$ satisfies these, so we just require a central square root preserving these

Thm (Rep cats of algebras) The following table gives the correspondence between algs + their rep cats:

Algebra A	\Rightarrow Cat $\text{Rep}(A)$
bialg	monoidal
Hopf alg	rigid monoidal
braided bialg	braided monoidal
braided Hopf alg	rigid braided monoidal
ribbon Hopf alg	ribbon

Proof Sketch

① bialg We saw last time that bialg $(A, m, \eta, \Delta, \epsilon)$ yields monoidal cat $(\text{Rep}(A), \otimes, F)$ as follows

- Given $V, W \in \text{ob}(\text{Rep}(A))$, $V \otimes W \in \text{ob}(\text{Rep}(A))$, via

$$a \cdot (v \otimes w) := \Delta(a) v \otimes w$$

- $1 \in \text{ob}(\text{Rep}(A))$ via

$$a \cdot 1 := \epsilon(a) 1$$

② Hopf alg We also saw that for Hopf alg $(A, m, \eta, \Delta, \epsilon, S)$, $\text{Rep}(A)$ is rigid w/ normal vec space duals + axn

$$(a \cdot v^*)(w) := v^*(S w)$$

\uparrow dual vec \uparrow vec

③ braided bialg Recall, for braided bialg $(A, m, \eta, \Delta, \epsilon, R)$, represented on V, W , we defined $c_{VW} = \tau_{VW} \circ R: V \otimes W \rightarrow W \otimes V$, + showed that these satisfy the Yang-Baxter eq. To prove that the c_{VW} are (components of) a braiding on $\text{Rep}(A)$, it remains to check that they are A -linear automorphisms:

$$\begin{aligned}
 c_{VW}(a \cdot (s \otimes t)) &= \tau_{VW}(R \Delta(a)(s \otimes t)) && \text{(defs of } c \text{ \& } A\text{-axn on } V \otimes W) \\
 &= \tau_{VW}(\Delta^{\text{op}}(a) R(s \otimes t)) && \text{(axiom of quasicoalgebra bialg)} \\
 &= \Delta(a) \tau_{VW}(R(s \otimes t)) && \text{(def of } \Delta^{\text{op}}) \\
 &= a \cdot (c_{VW}(s \otimes t)). && \text{(defs of } c \text{ \& } A\text{-axn on } V \otimes W)
 \end{aligned}$$

Finally, the hexagon identity follows from Yang-Baxter, together w/ the easily-verified identities,

$$\begin{aligned}
 c_{U, V \otimes W} &= (\text{id} \otimes c_{VW})(c_{UV} \otimes \text{id}) \\
 c_{U \otimes V, W} &= (c_{UV} \otimes \text{id})(\text{id} \otimes c_{VW})
 \end{aligned}$$

④ braided Hopf alg (combine ② + ③)

⑤ ribbon Hopf alg We need only equip $\text{Rep}(A)$ w/ a twist $\Theta_V: V \rightarrow V$; define $\Theta_V(w) = \Theta^{-1} \cdot w$.

Then Θ_V is an A -linear automorphism because $\Theta \in A$ is central + invertible. Now, we check the balancing axioms,

$$\begin{aligned}
& (\theta_V \otimes \theta_W) C_{uv} C_{vw} (s \otimes t) \quad \leftarrow s \in V, t \in W \\
&= (\theta^{-1} \otimes \theta^{-1}) (R_{21} R) (s \otimes t) \quad \text{(defs of twist } \theta \text{ \& braiding } c) \\
&= \Delta(\theta^{-1})(s \otimes t) \quad \text{(ribbon Hopf alg axiom)} \\
&= \theta_{V \otimes W} (s \otimes t), \quad \text{(def of } A\text{-axn on } V \otimes W)
\end{aligned}$$

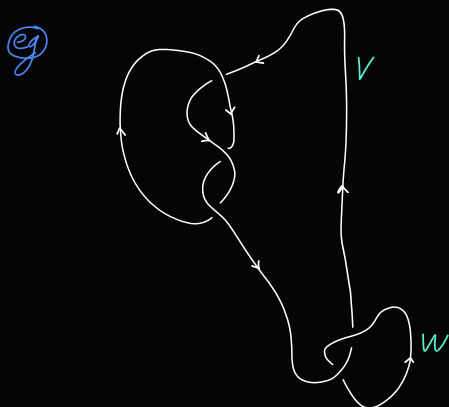
$$\begin{aligned}
& (\theta_V)^*(\alpha)(t) = \alpha(\theta_V(t)) \quad \leftarrow \alpha \in V^*, t \in V \quad \text{(def of dual map } (\theta_V)^*) \\
&= \alpha(\theta^{-1}t) \quad \text{(def of twist } \theta) \\
&= \alpha(S(\theta^{-1})t) \quad \text{(ribbon Hopf alg axiom)} \\
&= (\theta^{-1}\alpha)(t) \quad \text{(def of axn on } V^*) \\
&= [\theta_{V^*}(\alpha)](t) \quad \text{(def of twist } \theta) \quad \blacksquare
\end{aligned}$$

Def (Link coloring) Given a ribbon Hopf alg $(A, m, \eta, \Delta, \epsilon, S, R, \theta)$, and a link \mathcal{L} , a "coloring" of \mathcal{L} by reps is an assignment $\{L_1, \dots, L_n\} \rightarrow \{V_1, \dots, V_m\}$ of reps V_i of A to link components L_i of \mathcal{L} .

Denote a colored link $\mathbb{V}(\mathcal{L})$.

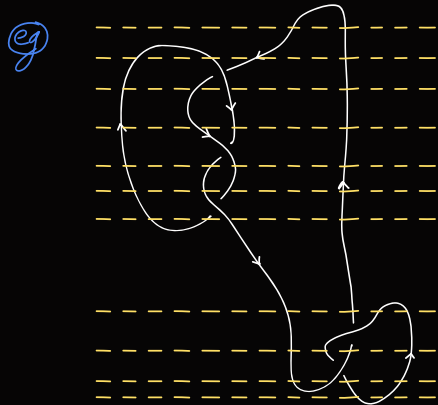
Worked Example (Intertwiner associated to colored link)

- ① Let $(A, m, \eta, \Delta, \epsilon, S, R, \theta)$ be a ribbon Hopf alg w/ S invertible; & $\mathbb{V}(\mathcal{L})$ be a colored oriented framed link
- ② Embed \mathcal{L} in $S^3 \cong \mathbb{R}^3 \cup \{\infty\}$; pick a plane in S^3 & isotopy \mathcal{L} to have blackboard framing wrt the plane



③ Interpret $V(L)$ as the string diagram encoding some A -linear morphism $F \rightarrow F$; explicitly,

④ isotopy $V(L)$ st its projection decomposes into vertically-stacked slices, each containing exactly 1 of the following



⑤ associate A -linear homomorphisms as follows:

	$c \equiv \tau \circ R$	
	c^{-1}	
	$1_F \mapsto \sum e^k \otimes u^{-1} \theta e_k$	$F \rightarrow V^* \otimes V$
	$1_F \mapsto \sum e^k \otimes e_k$	$F \rightarrow V \otimes V^*$
	$v \otimes \alpha \mapsto \alpha(\theta^{-1} u v)$	$V \otimes V^* \rightarrow F$
	$\alpha \otimes v \mapsto \alpha(v)$	$V^* \otimes V \rightarrow F$

recall, with $R = \sum \alpha_j \otimes \beta_j$
 $u = \sum s(\beta_j) \alpha_j$

remark we can check from 1st principles that all of the above is A -linear, but it is messy

eg for $v \otimes \alpha \mapsto \alpha(\theta^{-1} u v)$, we have

$$\begin{aligned}
 \phi(a \cdot (v \otimes \alpha)) &= \phi(\Delta(a)(v \otimes \alpha)) && \text{(def of } A\text{-axn on } V \otimes W) \\
 &= \phi\left(\sum_{(a')} a' v \otimes a' \alpha\right) && \text{(rearrangement)} \\
 &= \alpha\left(\sum_{(a')} s(a') \theta^{-1} u a' v\right) && \text{(defs of } \phi \text{ \& } A\text{-axn on } V^*) \\
 &= \alpha\left(\sum_{(a')} s(a') \theta^{-1} S^2(a') u v\right) && (S^2(-) = u(-)u^{-1}) \\
 &= \alpha\left(\sum_{(a')} s(a') S^2(a') \theta^{-1} u v\right) && (\theta^{-1} \in Z(A))
 \end{aligned}$$

$$\begin{aligned}
&= \alpha(S(m(\text{solid})) \Delta(a) \theta^{-1}uv) && \text{(rearrangement)} \\
&= \alpha(S(\varepsilon(a)1) \theta^{-1}uv) && \text{(def of } S) \\
&= \varepsilon(a) \alpha(\theta^{-1}uv) && \text{(linearity)} \\
&= a \cdot \phi(\alpha \otimes v) && \text{(defs of } A\text{-axn on } F)
\end{aligned}$$

④ The result of ③ is some intertwiner $\Psi: F \rightarrow F$ represented by $\mathbb{V}(\mathcal{L})$

Def (Reshetikhin-Turaev invariant) The Reshetikhin-Turaev invariant associated to an oriented framed link $\mathbb{V}(\mathcal{L})$, coloured by reps of some ribbon Hopf alg $(A, m, \eta, \Delta, \varepsilon, S, R, \theta)$ is the scalar $\langle \mathbb{V}(\mathcal{L}) \rangle \in F$ by which the morphism $\Psi: F \rightarrow F$ encoded by $\mathbb{V}(\mathcal{L})$ as above multiplies inputs

$$\text{ie. } \Psi: z \mapsto \langle \mathbb{V}(\mathcal{L}) \rangle z$$

Thm $\langle \mathbb{V}(\mathcal{L}) \rangle$ is a bona fide link invariant

Proof (special case of coherence thm for ribbon cats) ■

↑ 1st principles proof also found in Gelca (2014), thm 8.2


Thm (Identities for $\langle \mathbb{V}(\mathcal{L}) \rangle$) Denote $\mathbb{V}(L_i) = V_i$ for link comp L_i ; Then

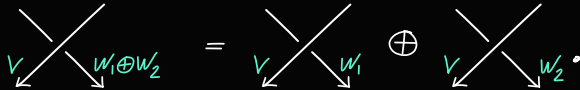
- ① If $V_j = W_1 \oplus W_2$, then $\langle \mathbb{V}(\mathcal{L}) \rangle = \langle \mathbb{V}'(\mathcal{L}) \rangle + \langle \mathbb{V}_2(\mathcal{L}) \rangle$, where

$$V_i(L_j) = \begin{cases} V_k & k \neq j \\ W_i & k = j \end{cases}$$
- ② (cabling principle) If $V_j = W_1 \otimes W_2$, then $\langle \mathbb{V}(\mathcal{L}) \rangle = \langle \mathbb{V}'(\mathcal{L}') \rangle$, where
 - \mathcal{L}' is \mathcal{L} , but with $L_j \rightsquigarrow L_j^{\parallel 2}$
 - \mathbb{V}' is \mathbb{V} but assigning W_1 & W_2 to the 2 comps of $L_j^{\parallel 2}$
- ③ $\langle \mathbb{V}(\mathcal{L}) \rangle = \langle \mathbb{V}'(\mathcal{L}') \rangle$, where
 - \mathcal{L}' is \mathcal{L} , but with some L_j reversed
 - \mathbb{V}' is \mathbb{V} but assigning V_j^* to L_j
- ④ If V_j is the trivial rep, then $\langle \mathbb{V}(\mathcal{L}) \rangle = \langle \mathbb{V}'(\mathcal{L}') \rangle$, where
 - \mathcal{L}' is \mathcal{L} , but w/ L_j deleted
 - \mathbb{V}' is \mathbb{V} restricted to \mathcal{L}'

Proof Sketch

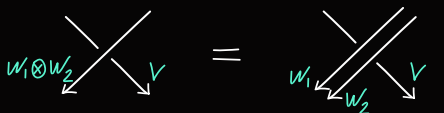
① Note that $a \circ (s \oplus t) = (as) \oplus (at)$, so the homomorphisms defined by $\times, \times, \cup, \cup$ split across \oplus . For instance, for a crossing

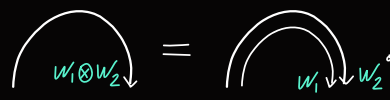
 we identify $V \otimes (W_1 \oplus W_2) = V \otimes W_1 \oplus V \otimes W_2$, whence the endomorphism defined by R acting on $V \otimes (W_1 \oplus W_2)$ splits into endomorphisms on $V \otimes W_1$ and $V \otimes W_2$; the same is then true for c & c^{-1} , so

 Similar reasoning gives splittings of the other homomorphisms.

② Consider first the homomorphism c associated to \times . We compute

$$\begin{aligned} c_{W_1 \otimes W_2, V} &= \tau_{W_1 \otimes W_2, V} R_{W_1 \otimes W_2, V} && \text{(def of } c) \\ &= \tau_{W_1 \otimes W_2, V} (\Delta \otimes id)(R)_{(W_1 \otimes W_2) \otimes V} && \text{(rearrangement)} \\ &= \tau_{W_1 \otimes W_2, V} (R_{13} R_{23})_{W_1 \otimes W_2 \otimes V} && \text{(axiom of braided bialg)} \\ &= \sum_{j,k} \beta_j \beta_k \otimes \alpha_j \otimes \alpha_k && \text{(writing } R = \sum \alpha_j \otimes \beta_j) \\ &= (c_{W_1, V} \otimes id)(id \otimes c_{W_2, V}) && \text{(def of } c) \end{aligned}$$

ie, 

We can similarly easily check that 

The case of $v \otimes \alpha \xrightarrow{\Phi} \alpha(\theta^{-1}uv)$ requires a bit more work + relies on the other braided bialg property; we omit the details (see Gela (2014), §8.1.4)

③ Reversing a strand + colouring w/ the dual encodes the same rep; the only "issues" are at maxima/minima, but our pivotal structure ensures consistency

④ Trivial rep $V^0 = F$; now $R: V \otimes F \rightarrow V \otimes F$ is the map

$$V \rightarrow V: v \mapsto \sum_j \varepsilon(\beta_j) \alpha_j v, \text{ by def of } A\text{-axn on } F.$$

$$\text{But } \sum_j \varepsilon(\beta_j) \alpha_j = m(\varepsilon \otimes id)(R) = 1 \text{ (thm from last time),}$$

$$\text{so } c: V \otimes F \rightarrow V \otimes F \equiv id_V.$$

Moreover, since $\varepsilon(\theta) = \varepsilon(u) = 1$, each of the other homomorphisms (cups + caps) also act as identity under identification $F \otimes F = F$ 